# Supplement to <br> "Wasteful sanctions, underperformance, and endogenous supervision" 

By David A. Miller and Kareen Rozen*


#### Abstract

This supplement studies optimality within a special class of contracts: those with increasingly harsh marginal penalties, which we term "decreasing convex" contracts. Within this class, simple contracts with work-target strategies and kinked-linear sanctioning schemes are optimal. JEL: C72, D03, D86 Keywords: Partnership, teams, moral hazard, monitoring, supervision, informal sanctions


## Optimal convex contracts

In this supplemental appendix we consider contracts that are symmetric with respect to task names and for which the amount of monitoring to be accomplished (denoted $F$ ) is public. In this case, the sanction depends on the number of failures $f$ of inspection, where $f \in\{0,1, \ldots, F\}$. Within this class, contracts which deliver increasingly large sanctions for larger numbers of inspection failures may be a focal class to consider. Such decreasing convex ( $D C$ ) contracts satisfy the restriction $v(f)-v(f+1) \geq v(f-1)-v(f) \geq 0$. Convex contracts may be natural in settings where sanctions are imposed by third parties who are more inclined to exact sanctions if they perceive a consistent pattern of failures. Conversely, a non-convex contract may be particularly difficult to enforce via an affected third party, since it would require leniency on the margin for relatively large injuries. For arbitrary capacity $M$, we show that DC contracts optimally induce work target strategies. Furthermore, the optimal such contract forgives failures up to some threshold, and increases the sanction linearly thereafter.

Theorem 5. For any $M$, work-target strategies with a kinked linear sanctioning scheme are optimal in the class of DC contracts.

We first prove several lemmas. The first provides a sufficient condition on a one-parameter family of probability distributions for the expectation of a concave function to be concave in the parameter. Though it can be derived as a corollary of a more general theorem of Susan Athey (2000), we provide a simple statement of

[^0]the condition along with a direct proof. We say that a function $\psi:\{0,1, \ldots, R\} \rightarrow$ $\mathbb{R}$ is concave if $\psi(r+1)-\psi(r) \leq \psi(r)-\psi(r-1)$ for all $r=1, \ldots, R-1$. A function $\phi: \mathcal{Z} \rightarrow \mathbb{R}$, where $\mathcal{Z} \subseteq \mathbb{R}$, is double crossing if there is a (possibly empty) convex set $A \subset \mathbb{R}$ such that $\bar{A} \cap \mathcal{Z}=\{z \in \mathcal{Z}: \phi(z)<0\}$.
Lemma 5 (Preservation of concavity). Let $\mathcal{R}=\{0,1, \ldots, R\}$, and let $\left\{q_{z}\right\}_{z \in \mathcal{Z}}$ be a collection of probability distributions on $\mathcal{R}$ parameterized by $z \in \mathcal{Z}=\{0,1, \ldots, Z\}{ }^{1}$ The function $\Psi(z)=\sum_{r=0}^{R} \psi(r) q_{z}(r)$ is concave if

1) There exists $k, c \in \mathbb{R}, k \neq 0$, such that $z=k \sum_{r=0}^{R} r q_{z}(r)+c$ for all $z \in \mathcal{Z}$;
2) $q_{z+1}(r)-2 q_{z}(r)+q_{z-1}(r)$ for all $z=1, \ldots, Z-1$, as a function of $r$, is double crossing;
3) $\psi:\{0,1, \ldots, R\} \rightarrow \mathbb{R}$ is concave.

Proof. Since $z=k \sum_{r=0}^{R} r q_{z}(r)+c$, there exists $\hat{b} \in \mathbb{R}$ such that $\sum_{r=0}^{R}(m r+$ b) $q_{z}(r)=\frac{m}{k} z+\hat{b}+c$ for any real $m$ and $b$. Hence, for any $m$ and $b$,

$$
\begin{equation*}
\sum_{r=0}^{R}(m r+b)\left(q_{z+1}(r)-2 q_{z}(r)+q_{z-1}(r)\right)=\frac{m}{k}(z+1-2 z+z-1)=0 \tag{B1}
\end{equation*}
$$

for all $z=1, \ldots, Z-1$. Therefore, for any $m$ and $b$, the second difference of $\Psi(z)$ is

$$
\begin{align*}
\Psi(z+1)-2 \Psi(z)+\Psi(z-1) & =\sum_{r=0}^{R} \psi(r)\left(q_{z+1}(r)-2 q_{z}(r)+q_{z-1}(r)\right)  \tag{B2}\\
& =\sum_{r=0}^{R}(\psi(r)-m r-b)\left(q_{z+1}(r)-2 q_{z}(r)+q_{z-1}(r)\right) .
\end{align*}
$$

By assumption, $q_{z+1}(r)-2 q_{z}(r)+q_{z-1}(r)$, as a function of $r$, is double crossing. Furthermore, since $\psi$ is concave, we can choose $m$ and $b$ such that, wherever $\left(q_{z+1}(r)-2 q_{z}(r)+q_{z-1}(r)\right)$ or $\frac{\partial^{2}}{\partial z^{2}} q_{z}(r)$ is nonzero, $\psi(r)-m r-b$ either has the opposite sign or is zero. From Eq. B2 we may conclude $\Psi(z)$ is concave.

The next lemma says that the expected sanctioning scheme will be decreasing convex in the number of tasks completed ${ }^{2}$
Lemma 6. If $v$ is decreasing convex, then $h_{v} \equiv \sum_{f=0}^{F} v(f) g(f, \cdot)$ is decreasing convex.

[^1]Proof. By letting $a \equiv|A|$, reversing the order of summation, and using fact that $\binom{k}{f}=0$ when $k<f$, we can write $h_{v}(A)$ as follows:

$$
\begin{align*}
h_{v}(A) & =\sum_{f=0}^{F} g(f, a) v(f) \\
& =\sum_{f=0}^{F}\left(\sum_{k=0}^{F} \frac{\binom{p-a}{k}\binom{a}{F-k}}{\binom{p}{F}}\binom{k}{f} \gamma^{f}(1-\gamma)^{k-f}\right) v(f)  \tag{B3}\\
& =\sum_{k=0}^{F} \frac{\binom{p-a}{k}\binom{a}{F-k}}{\binom{p}{F}}\left(\sum_{f=0}^{F}\binom{k}{f} \gamma^{f}(1-\gamma)^{k-f} v(f)\right) .
\end{align*}
$$

Therefore, the expectation is first with respect to the binomial, and then with respect to the hypergeometric. Applying Lemma 5 twice gives the result. First, note that the expectation of the binomial is $\gamma k$, a linear function of $k$, while the expectation of the hypergeometric is $\frac{F}{p}(p-a)$, a linear function of $a$. Hence it suffices to show that the binomial second-difference in $k$ is double-crossing in $f$ (hence the inside expectation is decreasing convex in $k$ ) and the hypergeometric second-difference in $a$ is double-crossing in $k$. To see this is true for the binomial, note that we may write the binomial second-difference in $k$ as

$$
\begin{equation*}
\binom{k}{f} \gamma^{f}(1-\gamma)^{k-f}\left(\frac{(k+1)(1-\gamma)}{k+1-f}-2+\frac{k-f}{k(1-\gamma)}\right) . \tag{B4}
\end{equation*}
$$

It can be shown that the term in parentheses is strictly convex in $f$ and therefore double crossing in $f$, so the whole expression is double-crossing in $f$. To see this is true for the hypergeometric, note that we may write the hypergeometric second-difference in $a$ as

$$
\begin{equation*}
\frac{\binom{p-a}{k}\binom{a}{F-k}}{\binom{p}{F}}\left(\frac{p-a-k}{p-a} \cdot \frac{a+1}{a+1-F+k}-2+\frac{p-a+1}{p-a+1-k} \cdot \frac{a-F+k}{a}\right) . \tag{B5}
\end{equation*}
$$

It can be shown that the term in parentheses has either zero or two real roots $3^{3}$ If there are no real roots, then the term in parentheses is double-crossing in $k$ (the region in which it is negative must be convex, but may be empty), and thus the whole expression is double-crossing in $k$. If there are two real roots, it can be shown that the derivative with respect to $k$ is negative at the smaller root, and thus both the term in parentheses and the whole expression are double-crossing in $k$.

[^2]Proof of Theorem 5. Fix any $p, F$, and $\lambda$. Suppose strategy $s$, with $p^{*}>0$ the maximal number of tasks completed, is optimal. Consider the decreasing convex contract $v$ that implements $s$ at minimum cost. Because $v$ is decreasing, MLRP (or FOSD in $a$ ) implies the expected sanction decreases in the number of completed tasks: $h(a)>h(a-1)$ for all $a$. By contradiction, suppose the downward constraint for $p^{*}$ versus $p^{*}-1$ is slack: $h\left(p^{*}\right)-h\left(p^{*}-1\right)>c-b$. By Lemma 6 and monotonicity, for any $k>1, h\left(p^{*}-k+1\right)-h\left(p^{*}-k\right)>c-b$. But then for any $a$ with $s(a)=a$ and every $a^{\prime}<a$, the downward constraint $h(a)-h\left(a^{\prime}\right)=\sum_{k=a^{\prime}}^{a-1} h(k+1)-h(k) \geq\left(a-a^{\prime}\right)(c-b)$ is slack. Some constraint must bind at the optimum, else the strategy is implementable for free, so the downward constraint for $p^{*}$ versus $p^{*}-1$ must bind. Again, each downward constraint is satisfied, and for any $a>p^{*}, h(a)-h\left(p^{*}\right)<\left(a-p^{*}\right)(c-b)$. So the strategy $s$ has a work target of $p^{*}$.

Suppose we look for the optimal convex contract with $p$ assigned tasks, $F$ monitoring slots, and strategy $s$ with work target $p^{*}$. By the above, the only binding incentive constraint is the downward constraint for completing $p^{*}$ tasks. Since $v(0)=0$, convexity implies monotonicity. The constraint $v(0) \geq 0$ does not bind $\sqrt{4}^{4}$ so the cost minimization problem in primal form is

$$
\begin{align*}
& \max _{(-v) \geq 0} \sum_{f=0}^{F}\left(-(-v(f)) \sum_{a=0}^{p}-g(f, a) t_{s}(a)\right) \text { subject to } \\
& \quad \sum_{f=0}^{F}(-v(f))\left(g\left(f, p^{*}\right)-g\left(f, p^{*}-1\right)\right) \leq-(c-b),  \tag{B6}\\
& \quad 2(-v(f))-(-v(f+1))-(-v(f-1)) \leq 0 \text { for all } f=1, \ldots, F-1,
\end{align*}
$$

where $t_{s}(a)=\sum_{a^{\prime}=a}^{p} \mathbb{I}\left(s\left(a^{\prime}\right)=a\right)\binom{p}{a^{\prime}} \lambda^{a^{\prime}}(1-\lambda)^{p-a^{\prime}}$ is the probability of completing $a$ tasks given strategy $s$. Let $x$ be the Lagrange multiplier for the incentive compatibility constraint, $z_{f}$ the multiplier for the convexity constraint $2(-v(f))-$ $(-v(f+1))-(-v(f-1)) \leq 0$, and $\vec{z}$ the vector $\left(z_{1}, \ldots, z_{F-1}\right)$. The constraint set can be written $A^{\top} \cdot(-v(0), \ldots,-v(F))$, where, in sparse form,

$$
A=\left(\begin{array}{crcr}
g\left(0, p^{*}\right)-g\left(0, p^{*}-1\right) & -1 & &  \tag{B7}\\
\vdots & 2 & \ddots & \\
\vdots & -1 & \ddots & -1 \\
\vdots & & \ddots & 2 \\
g\left(F, p^{*}\right)-g\left(F, p^{*}-1\right) & & & -1
\end{array}\right)
$$

${ }^{4}$ Although $v(0) \geq 0$ is satisfied with equality, the binding constraint on $v(0)$ is actually $v(0) \leq 0$.

Let $r$ be the vector of dual variables: $r=\left(x, z_{1}, \ldots, z_{F-1}\right)$. The dual problem is

$$
\begin{equation*}
\min _{r \geq \overrightarrow{0}}(b-c) x \text { s.t. }(A r)_{f} \geq-\sum_{a=0}^{p} g(f, a) t_{s}(a) \text { for all } f=0,1, \ldots, F \tag{B8}
\end{equation*}
$$

where $(A r)_{f}$ is the $(f)$ th component of $A \cdot r$; i.e.,

$$
\begin{equation*}
(A r)_{f}=x\left(g\left(f, p^{*}\right)-g\left(f, p^{*}-1\right)\right)-z_{f-1}+2 z_{f}-z_{f+1} \tag{B9}
\end{equation*}
$$

where we define $z_{0} \equiv 0, z_{F} \equiv 0$, and $z_{F+1} \equiv 0$. Let $\hat{f}$ be the smallest $f$ with $v(f)<0$. It must be that $v(f)<0$ for all $f \geq \hat{f}$, so by duality, $(A \cdot r)_{f} \geq$ $-\sum_{a=0}^{p} g(f, a) t_{s}(a)$ binds for all $f \geq \hat{f}$. Hence

$$
\begin{equation*}
x=\frac{\sum_{a=0}^{p} g(f, a) t_{s}(a)-z_{f-1}+2 z_{f}-z_{f+1}}{g\left(f, p^{*}-1\right)-g\left(f, p^{*}\right)} \text { for all } f=\hat{f}, \ldots, F . \tag{B10}
\end{equation*}
$$

In particular, this means that if $z_{F-1}=0$ (implied for $\hat{f}=F$ ) then the optimal contract (which would have expected sanction $-x(c-b)$ ) has the same value as that derived in ??, completing the claim. Henceforth we assume $z_{F-1}>0$. The sum of the $z$-terms over $(A \cdot r)_{-1}$ and $(A \cdot r)_{F}$ is $-z_{F-1}+\left(2 z_{F-1}-z_{F-2}\right)=$ $z_{F-1}-z_{F-2}$. Note also the corresponding sum of $z$-terms over $F-2, F-1$, and $F:-z_{F-1}+\left(2 z_{F-1}-z_{F-2}\right)+\left(-z_{F-3}+2 z_{F-2}-z_{F-1}\right)=z_{F-2}-z_{F-3}$. Iterating, the sum of the $z$-terms in $(A \cdot r)_{f}$ from any $\tilde{f} \geq \hat{f}$ to $F$ is $z_{\tilde{f}}-z_{\tilde{f}-1}$. Summing the equalities in Eq. B10 thus yields a recursive system for $z_{\tilde{f}}$ for all $\tilde{f}=\hat{f}, \ldots, F$ :

$$
\begin{equation*}
z_{\tilde{f}}=z_{\tilde{f}-1}-\sum_{f=\tilde{f}}^{F} \sum_{a=0}^{p} g(f, a) t_{s}(a)+x \sum_{f=\tilde{f}}^{F}\left(g\left(f, p^{*}-1\right)-g\left(f, p^{*}\right)\right) . \tag{B11}
\end{equation*}
$$

By definition, the convexity constraint is slack at $\hat{f}-1$, so $z_{\hat{f}-1}=0$. By induction, for $f^{\prime}=\hat{f}, \ldots, F$,

$$
\begin{equation*}
z_{f^{\prime}}=-\sum_{\tilde{f}=\hat{f}}^{f^{\prime}} \sum_{f=\tilde{f}}^{F} \sum_{a=0}^{p} g(f, a) t_{s}(a)+x \sum_{\tilde{f}=\hat{f}}^{f^{\prime}} \sum_{f=\tilde{f}}^{F}\left(g\left(f, p^{*}-1\right)-g\left(f, p^{*}\right)\right) \tag{B12}
\end{equation*}
$$

Plugging Eq. B12 for $f^{\prime}=F$ into the binding constraint $(A r)_{F} \geq-\sum_{a=0}^{p} g(F, a) t_{s}(a)$
yields:

$$
\begin{equation*}
x=\frac{\sum_{\tilde{f}=\hat{f}}^{F} \sum_{f=\tilde{f}}^{F} \sum_{a=0}^{p} g(f, a) t_{s}(a)}{\sum_{\tilde{f}=\hat{f}}^{F} \sum_{f=\tilde{f}}^{F}\left(g\left(f, p^{*}-1\right)-g\left(f, p^{*}\right)\right)} . \tag{B13}
\end{equation*}
$$

The expectation of a random variable $X$ on $\{0, \ldots, n\}$, is $\sum_{j=1}^{n} j \operatorname{Pr}(X=j)$, which also equals $\sum_{j=1}^{n} \operatorname{Pr}(X \geq j)$. Since $\sum_{f=\tilde{f}}^{F} \sum_{a=0}^{p} g(f, a) t_{s}(a)=\operatorname{Pr}(f \geq \tilde{f})$, the numerator of Eq. B13 equals

$$
\begin{array}{r}
\sum_{\tilde{f}=\hat{f}}^{F} \sum_{f=\tilde{f}}^{F} \sum_{a=0}^{p} g(f, a) t_{s}(a)=\sum_{\tilde{f}=\hat{f}}^{F} \operatorname{Pr}(f \geq \tilde{f})=\sum_{\tilde{f}=\hat{f}}^{F}(\tilde{f}-\hat{f}+1) \operatorname{Pr}(f=\tilde{f})  \tag{B14}\\
=\sum_{\tilde{f}=1}^{F}(\tilde{f}-\hat{f}+1)_{+} \operatorname{Pr}(f=\tilde{f})=\mathbb{E}\left((f-\hat{f}+1)_{+}\right) \equiv \mathbb{E}(\phi(\hat{f}))
\end{array}
$$

where $(y)_{+} \equiv \max \{y, 0\}$ and $\phi$ is the random function $\phi(\hat{f}) \equiv(f-\hat{f}+1)_{+}$. In words, $\phi(\hat{f})$ is the number of discovered unfulfilled tasks that exceed the threshold for sanctions $\hat{f}$. The denominator of Eq. B13 can be rewritten similarly, yielding

$$
\begin{equation*}
x=\frac{\mathbb{E}(\phi(\hat{f}))}{\mathbb{E}\left(\phi(\hat{f}) \mid a=p^{*}-1\right)-\mathbb{E}\left(\phi(\hat{f}) \mid a=p^{*}\right)} . \tag{B15}
\end{equation*}
$$

The minimized expected sanction is $\mathbb{E}(v(f))=(b-c) x$, and is implemented by

$$
v(f)=-\frac{(c-b)(f-\hat{f}+1)_{+}}{\mathbb{E}\left(\phi(\hat{f}) \mid a=p^{*}-1\right)-\mathbb{E}\left(\phi(\hat{f}) \mid a=p^{*}\right)} \text { for all } f=0,1, \ldots, F
$$

## REFERENCES

Athey, Susan. 2000. "Characterizing properties of stochastic objective functions." Working paper.


[^0]:    * Miller: University of Michigan, Dept. of Economics, 611 Tappan St., Ann Arbor, MI 48109, econdm@umich.edu. Rozen: Yale, Dept. of Economics and the Cowles Foundation for Research in Economics, 30 Hillhouse Ave., New Haven, CT 06511, kareen.rozen@yale.edu.

[^1]:    ${ }^{1}$ A similar result holds if $z \in \mathcal{Z}=[0,1]$.
    ${ }^{2}$ Recall that $g(f, a) \equiv \sum_{k=f}^{F_{i}} \frac{\binom{p_{i}-a}{k}\binom{a}{F_{i}-k}}{\binom{p_{i}}{F_{i}}}\binom{k}{f} \gamma^{f}(1-\gamma)^{k-f}$.

[^2]:    ${ }^{3}$ The term in parentheses does not account for the fact that the entire expression equals zero whenever $k>p-a$ or $F-k>a$. However, on the closure of these regions the second difference cannot be negative, and so these regions may be ignored.

